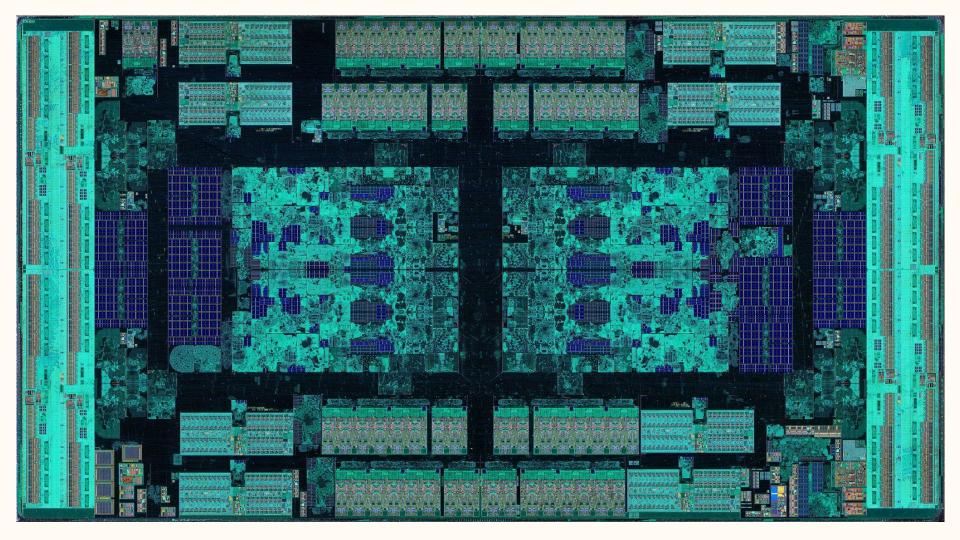
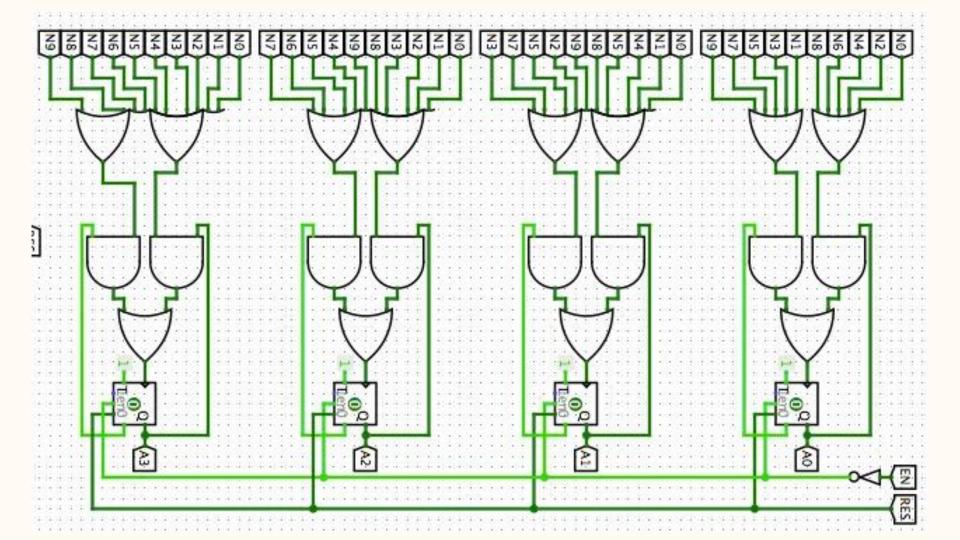


Assembly language: Where you write 10 lines of code to do what one line of Python can... but hey, at least you're closer to the machine!



0000000000000	01001000	11000111	11000111	Η
00000003:	00110111	00010011	000000000	7
00000006:	000000000	01001000	11000111	.Η.
00000009:	11000000	00111100	000000000	.<.
0000000c:	000000000	000000000	00001111	• • •
:+0000000	00000101			•



#### 00000000: 48 c7 c7 H... 00000003: 37 13 00 7.. 00000006: 00 48 c7 .H. 00000009: **C0** 3C 00 0000000c: 00 00 0f0000000f: 05

## 0000000048C7C737130000mov rdi,0x13370000000748C7C03C000000mov rax,0x3c0000000E0F05syscall

ARM AArch64

## mov x8, 93 mov x0, 0x1337 svc #0



## li \$v0, 4001 li \$a0, 0x1337 syscall

# mov eax, 1 mov ebx, 0x1337 int 0x80

## ld a, 0x37 ld b, 0x13 call 0x05



## li a7, 93 li a0, 0x1337 ecall

li r0, 1 li r3, 0x1337

SC

PPC

#### Registers

- 1. General-Purpose Registers (16 registers)
- 64-bit Registers:
  - RAX, RBX, RCX, RDX (traditional general-purpose registers)
  - RSI, RDI (used for string operations or function arguments)
  - RBP, RSP (base and stack pointers)
  - R8 to R15 (additional general-purpose registers in x86\_64)
- These registers can also be accessed in smaller chunks:
  - **32-bit**: EAX, EBX, ECX, EDX, etc.
  - **16-bit**: AX, BX, CX, DX, etc.
  - **8-bit**: AL, BL, CL, DL, etc.

#### 2. Special-Purpose Registers (6 primary ones)

- Instruction Pointer:
  - **RIP**: Holds the address of the next instruction to execute.
- Flags Register:
  - RFLAGS: Stores flags for arithmetic operations, control flow, etc.
- Segment Registers:
  - CS, DS, SS, ES, FS, GS: Mostly legacy, but FS and GS are still used in modern x86\_64 for things like thread-local storage.

#### 3. Floating-Point and Vector Registers (32 registers)

- XMM Registers (128-bit):
  - XMM0 to XMM15: Used for SIMD (Single Instruction, Multiple Data) operations.
- YMM Registers (256-bit):
  - YMM0 to YMM15: Used with AVX (Advanced Vector Extensions).
- ZMM Registers (512-bit):
  - ZMM0 to ZMM31: Available on processors with AVX-512 support.
- FPU Registers:
  - ST0 to ST7: Legacy floating-point registers from the x87 FPU stack.

#### 4. Control and Debug Registers

- Control Registers (4 primary ones):
  - CR0, CR2, CR3, CR4: Used for system-level settings like memory management.
- Debug Registers (8 registers):
  - DR0 to DR7: Used for setting hardware breakpoints and debugging.

- **5. Other Specialized Registers** 
  - Model-Specific Registers (MSRs): Configuration and performance monitoring.
  - **Test Registers** (legacy): Rarely used today.
  - **Performance Counters**: Used for profiling and optimization.

4	32	16 8	0	64
%rax	%eax	%ax 8ah	tal	%r8
%rbx	%ebx	%bx %bh	<b>b</b> bl	%r9
%rc <b>x</b>	%ecx	*cx *ch	łcl	% <b>r10</b>
%rdx	%edx	۶dx ۶dh	8d1	%r11
%rsp	fesp	۶sp ا	spl	%r12
%rbp	*ebp	%bp %	bpl	%r13
%rsi	%esi	%si <b>%</b>	sil	% <b>r14</b>
%rdi	%edi	۶di ۶	dil	%r15



#### Operations

#### **1. Arithmetic Operations**

- **Description**: Perform basic mathematical computations.
- Examples:
  - ADD, SUB Addition and subtraction.
  - MUL, IMUL Unsigned and signed multiplication.
  - DIV, IDIV Unsigned and signed division.
  - INC, DEC Increment and decrement.
  - ADC, SBB Add and subtract with carry/borrow.

#### 2. Logical Operations

- **Description**: Perform bitwise and logical computations.
- Examples:
  - AND, OR, XOR Bitwise AND, OR, and XOR.
  - NOT Bitwise negation.
  - TEST Perform a bitwise AND and set flags without storing the result.
  - CMP Compare two values by subtracting and setting flags.

#### **3. Data Movement Instructions**

- **Description**: Transfer data between registers, memory, and I/O.
- Examples:
  - MOV Move data between registers and memory.
  - PUSH, POP Push and pop values onto/from the stack.
  - LEA Load the effective address of a memory operand.
  - XCHG Exchange the contents of two locations.
  - CMOVcc Conditional move based on flags (e.g., CMOVE, CMOVNE).

#### **4. Control Flow Instructions**

- **Description**: Alter the flow of execution.
- Examples:
  - JMP Unconditional jump.
  - JE, JNE, JG, JL, etc. Conditional jumps based on flags.
  - CALL, RET Call a procedure and return from it.
  - LOOP Loop with a counter.

#### 5. String and Memory Operations

- **Description**: Operate on strings and memory blocks efficiently.
- Examples:
  - MOVSB, MOVSW, MOVSD Move string data.
  - STOSB, STOSW, STOSD Store string data.
  - LODSB, LODSW, LODSD Load string data.
  - CMPSB, CMPSW, CMPSD Compare string data.
  - SCASB, SCASW, SCASD Scan string data.

#### 6. Shift and Rotate Instructions

- **Description**: Shift and rotate bits in registers or memory.
- Examples:
  - SHL, SHR Shift left and right logically.
  - SAR Shift right arithmetically.
  - ROL, ROR Rotate bits left and right.
  - RCL, RCR Rotate bits through the carry flag.

#### 7. Input/Output Instructions

- **Description**: Read from or write to I/O ports.
- Examples:
  - IN, OUT Read from and write to an I/O port.
  - INSB, INSW, INSD Input from port to string.
  - OUTSB, OUTSW, OUTSD Output string to port.

#### 8. Floating-Point and SIMD Instructions

- **Description**: Perform floating-point arithmetic and vectorized operations.
- Examples:
  - FADD, FSUB, FMUL, FDIV Floating-point arithmetic.
  - MOVAPS, ADDPS, MULPS SIMD operations with packed single-precision floats.
  - PADDQ, PSLLD Integer SIMD operations.
  - SQRTPS, MINPS Specialized SIMD instructions.

#### 9. System-Level Instructions

- **Description**: Manage processor state, system calls, and privileged operations.
- Examples:
  - SYSCALL, SYSRET System call and return (Linux and Windows).
  - CPUID Get processor information.
  - HLT Halt the processor.
  - $\circ$  INT n Trigger a software interrupt.
  - IRET Return from an interrupt handler.

#### **10. Miscellaneous Instructions**

- **Description**: Instructions that don't fit cleanly into other categories.
- Examples:
  - NOP No operation.
  - PAUSE Hint to the CPU to reduce power or delay.
  - XLAT Translate a byte using a lookup table.
  - UD2 Undefined instruction (for debugging purposes).

#### syscall

#### What is syscall?

- The **syscall** instruction is used to make **system calls** in x86\_64 architecture.
- It transitions control from user mode to kernel mode, allowing programs to request services from the operating system (e.g., file I/O, process management).

#### How syscall Works

#### 1. Registers Used for Arguments:

- System call number: RAX
- Arguments:
  - RDI: First argument
  - RSI: Second argument
  - RDX: Third argument
  - R10: Fourth argument
  - **R8**: Fifth argument
  - R9: Sixth argument

#### How syscall Works

- 2. Registers Affected:
  - Return value: Stored in RAX after the syscall.
  - Flags: RFLAGS may change based on syscall results.

#### How syscall Works

#### 3. Instruction Flow:

- Load the syscall number into RAX.
- Load any required arguments into the appropriate registers.
- Execute syscall.
- Check the return value in RAX.

%rax	System call	%rdi	%rsi	%rdx	%r10	%r8	%r9
0	sys_read	unsigned int fd	char *buf	size_t count			
1	sys_write	unsigned int fd	const char *buf	size_t count			
2	sys_open	const char *filename	int flags	int mode			
3	sys_close	unsigned int fd					
4	sys_stat	const char *filename	struct stat *statbuf				
5	sys_fstat	unsigned int fd	struct stat *statbuf				
6	sys_lstat	fconst char *filename	struct stat *statbuf				
7	sys_poll	struct poll_fd *ufds	unsigned int nfds	long timeout_msecs			
8	sys_lseek	unsigned int fd	off_t offset	unsigned int origin			
9	sys_mmap	unsigned long addr	unsigned long len	unsigned long prot	unsigned long flags	unsigned long fd	unsigned long off
10	sys_mprotect	unsigned long start	size_t len	unsigned long prot			

## What is the call Instruction?

- The **call** instruction is used to invoke a **subroutine** (function).
- It performs two key tasks:
  - 1. **Pushes the return address** (the address of the next instruction after call) onto the stack.
  - 2. **Transfers control** to the subroutine by jumping to the specified address.

# How call Works

# 1. Push Return Address:

- The address of the instruction immediately after the call is pushed onto the stack.
- This ensures the program knows where to return after the subroutine finishes.
- 2. Jump to Target Address:
  - Control is transferred to the subroutine by jumping to the target address or label.

# Key Points:

# 1. Pairs with ret:

- The subroutine uses ret to pop the return address from the stack and jump back.
- 2. Indirect Calls:
  - You can use call with a register or memory address for indirect subroutine calls:

Ο

```
mov rax, my_function
call rax
```

# Writing, Compiling

> vi input.s

```
.intel_syntax noprefix
.globl start
```

\_start: mov rax, 60 mov rdi, 1337 syscall

# 1a. assemble

• as -o output.o input.s

## What is an Object File?

- An intermediate binary file produced by the assembler (e.g., as in GNU toolchain).
- Contains **machine code** and metadata required for linking and creating an executable.

#### Key Sections in an Object File

Section	Purpose	
.text	Contains the machine code (instructions) generated from the assembly source.	
.data	Stores initialized global and static variables (e.g., int $x = 42$ ; ).	
.bss	Holds uninitialized variables, which are zero-initialized at runtime (e.g., int y; ).	
.rodata	Contains read-only data such as string literals (e.g., "Hello, World!").	
.symtab	Symbol table, listing functions and global variables with associated metadata (names, addresses, etc.).	
.rel.text	Relocation table for .text , indicating places where addresses need adjustment during linking.	
.debug	Debugging information (optional, if compiled with debug flags).	

# **Common Tools to Inspect Object Files**

- 1. **objdump**: Disassembles and analyzes the object file.
  - Example: objdump -d input.o (disassembles the .text section).
- 2. **readelf**: Displays the object file structure.
  - Example: readelf -a input.o (shows all sections, symbols, and relocation info).

## Takeaways

- **Object files are not executables** but are crucial for the linking stage.
- They combine **code**, **data**, **and metadata** to facilitate building the final binary.

# 1b. Manual linking

• Id -o my-elf output.o

# **Definition:**

- Linking is the process of combining **object files** and **libraries** into a single executable binary.
- It resolves references between symbols (e.g., functions, variables) defined in different files.

# **Two Types of Linking:**

- 1. Static Linking:
  - Libraries are directly embedded into the executable.
  - Produces a standalone binary but increases size.

# 2. Dynamic Linking:

- External libraries are loaded at runtime.
- Reduces binary size but depends on system-installed libraries.

# Why Linking Matters?

- Combines code from multiple sources.
- Resolves function and variable dependencies.
- Optimizes and prepares a binary for execution.

# Key Steps Performed by 1d:

- 1. Symbol Resolution:
  - Matches undefined symbols (e.g., printf) to their definitions in libraries or other object files.
- 2. **Relocation**:
  - Adjusts memory addresses for symbols and code to match the final binary layout.
- 3. Section Merging:
  - Combines similar sections (e.g., .text, .data) from different object files.
- 4. Library Linking:
  - Includes required library functions based on symbol usage.

# 1c. Pull out code

• objcopy --dump-section .text=code my-elf

# 2. Assemble and Link one step

• gcc -nostdlib -static -o my-elf input.s

- 3. Compile straight to bytes
- > vi input.s

#### BITS 64

- start: mov rdi, 1337 mov rax, 60 syscall
- nasm -f bin input.s

# 4. pwntools

from pwn import asm, context

```
# Set architecture (x86_64 for 64-bit systems)
context.arch = 'amd64'
```

#### # Assembly instructions

assembly	_code = '	
mov rax,	60	; syscall: exit
mov rdi,	0	; exit code: 0
syscall		; make the syscall

```
# Assemble the code
machine_code = asm(assembly_code)
```

# Print the resulting machine code in hex format
print("Assembly Code:")
print(assembly\_code)

```
print("\nMachine Code:")
print(machine_code.hex())
```

# Running and Debugging

- **gdb**: General-purpose debugger for assembly and other languages.
- **pwntools**: Python library with built-in debugging utilities.
- **strace**: Traces system calls for insight into program behavior.

## What is gdb?

- The GNU Debugger (gdb) allows you to:
  - Step through assembly instructions.
  - Inspect registers and memory.
  - Set breakpoints to pause execution.

#### Key Commands:

Command	Description
gdb <program></program>	Start debugging a program.
break <addr></addr>	Set a breakpoint at an address or function.
run	Run the program until a breakpoint.
stepi Or si	Execute the next assembly instruction.
info registers	Display all register values.
x/ <n> <addr></addr></n>	Examine memory at an address ( <n> bytes).</n>

gdb program

- > break \*0x401000
- > stepi
- > info registers
- > x/10xw \$rsp

- # Set a breakpoint at an address
- # Run the program
- # Step through instructions
- # Check register state
- # Examine 10 words at RSP

## What is pwntools?

- Python library for binary exploitation and debugging.
- Provides tools for **dynamic debugging** using scripts.

#### What is strace?

- A tool to trace **system calls** made by a program.
- Helps debug issues related to:
  - File I/O.
  - Memory allocation.
  - Permissions or resource errors.

#### The Stack

## **Definition:**

- The **stack** is a region of memory used for temporary storage in programs.
- It operates in a Last In, First Out (LIFO) manner.

# **Key Characteristics:**

# 1. **Dynamic Allocation**:

- Automatically allocates and deallocates memory during function calls.
- 2. Directional Growth:
  - On x86\_64, the stack grows **downward** (toward lower memory addresses).
- 3. Managed by Registers:
  - **RSP**: Stack Pointer (points to the top of the stack).
  - **RBP**: Base Pointer (used for referencing local variables).

#### Why Use the Stack?

- Function Calls: Store return addresses, arguments, and local variables.
- **Temporary Storage**: Efficient for short-lived data.
- **Control Flow**: Helps manage recursive and nested functions.

#### **Basic Operations:**

- 1. **Push**: Adds data to the top of the stack.
  - Decreases RSP.

```
push rax ; Store RAX on the stack
```

- 2. Pop: Removes data from the top of the stack.
  - Increases RSP.

pop rax ; Restore the top value into RAX

# **Function Call Example:**

- 1. Caller pushes arguments onto the stack.
- 2. The return address is pushed automatically during call.
- 3. The callee allocates space for local variables.

High Memory Addresses

 Arguments
 <- Caller pushes arguments</td>

 Return Address
 <- CALL instruction pushes return address</td>

 Local Variables

 ------/
 <- RSP (Stack Pointer)</td>

Low Memory Addresses

# Calling Conventions:

- Defines how arguments, return values, and stack management are handled.
- x86\_64 Linux (System V ABI):
  - **Registers**: First 6 arguments in RDI, RSI, RDX, RCX, R8, R9.
  - **Stack**: Additional arguments and return address.

**Prologue (Callee Setup):** 

Save the previous base pointer: push rbp mov rbp, rsp

Allocate space for local variables:

sub rsp, <size>

#### **Epilogue (Callee Cleanup):**

Deallocate local variables:

add rsp, <size>

Restore the base pointer and return:

pop rbp

ret

\_start:

```
; Push values onto the stack
mov rax, 42
push rax ; Store 42 on the stack
mov rax, 100
push rax ; Store 100 on the stack
```

```
; Pop values off the stack
pop rbx ; RAX = 100
pop rcx ; RCX = 42
```

```
; Exit syscall
mov rax, 60
xor rdi, rdi
syscall
```



#### > challenges

### 1. Hello, World! (Data Movement + String Operations)

- **Goal**: Print "Hello, World!" to the screen using a syscall.
- Instructions: mov, syscall.
- Hints:
  - Use the write syscall (rax = 1) with the string in memory.
  - Pass the file descriptor (stdout = 1), string pointer, and length.

## 2. Add Two Numbers (Arithmetic Operations)

- **Goal**: Prompt the user to input two numbers, add them, and print the result.
- Instructions: add, mov, syscall.
- Hints:
  - Use the read syscall to get input.
  - Convert ASCII input to integers and use add.

# 3. Compare Two Numbers (Control Flow)

- **Goal**: Compare two user-provided numbers and print which one is larger.
- Instructions: cmp, jne, jl, jmp.
- Hints:
  - Use cmp to compare values and conditional jumps (j1, jg) to handle output.

### 4. Implement a Simple Loop (Control Flow + Arithmetic)

- **Goal**: Print numbers 1 through 10 in a loop.
- Instructions: mov, add, cmp, jmp.
- Hints:
  - Use a counter in a register.
  - Use cmp and jmp to create a loop.

### 5. Bitwise Manipulation (Logical Operations)

- **Goal**: Toggle the case of a string (convert uppercase to lowercase and vice versa).
- Instructions: xor, and, or.
- Hints:
  - Use bitwise xor with  $0 \times 20$  to toggle case.
  - Loop through each character in the string.

### 6. Shift and Rotate (Bitwise Operations)

- **Goal**: Multiply a number by 16 using a left shift.
- Instructions: shl, sar, mov.
- Hints:
  - Use sh1 to shift bits to the left.
  - Print the result using the write syscall.

# 7. Basic String Reverse (String and Memory Operations)

- **Goal**: Reverse a user-provided string.
- Instructions: movsb, rep, jmp.
- Hints:
  - Use pointers to swap characters in memory.
  - Iterate until the midpoint of the string.

# 8. Implement an XOR Cipher (Logical Operations)

- **Goal**: Encrypt a string using an XOR cipher with a fixed key.
- Instructions: xor, mov, loop.
- Hints:
  - XOR each character with a key (e.g.,  $0 \times AA$ ).
  - Print the encrypted result.

#### 9. Smallest Number in an Array (Arithmetic + Loops)

- Goal: Find the smallest number in an array of integers.
- Instructions: cmp, mov, jmp.
- Hints:
  - Use a register to store the smallest number.
  - Iterate through the array with a loop, updating the register when a smaller number is found.

### 10. Fibonacci Sequence (Advanced Control Flow + Arithmetic)

- **Goal**: Compute and print the first 10 numbers in the Fibonacci sequence.
- Instructions: mov, add, push, pop, jmp.
- Hints:
  - Use two registers to store the last two Fibonacci numbers.
  - Loop to calculate and print each new number.